Hecke insertion and maximal increasing and decreasing sequences in fillings of polyominoes

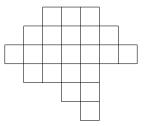
School of Mathematical and Statistical Sciences Clemson University

joint work with Ting Guo

Triangle Lectures in Combinatorics University of North Carolina at Charlotte

February 29, 2020

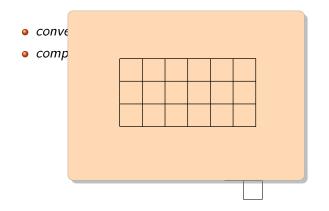
- convex rows and columns
- comparable rows and columns



• Lengths of rows from top to bottom form a unimodal sequence.

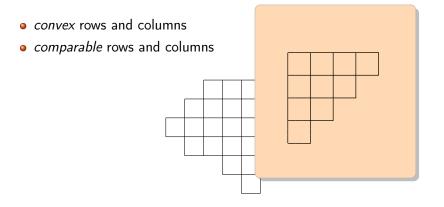
(Clemson)

Moon Polyominoes

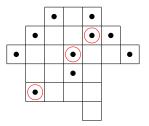


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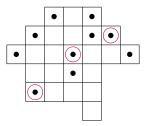
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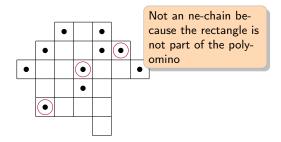
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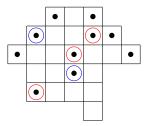
- ne-chain: A set of 1-cells {(i₁, j₁), (i₂, j₂), ..., (i_k, j_k)} with i₁ < ··· < i_k, j₁ < ··· < j_k such that the smallest rectangle containing them is a subset of the polyomino.
- ne(M) = the length of the largest ne-chain in the filling M.



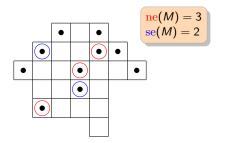
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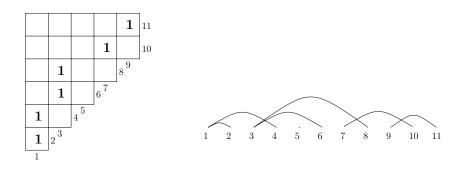


- se-chain: A set of 1-cells {(i₁, j₁), (i₂, j₂), ..., (ik, jk)} with i₁ < ··· < ik, j₁ > ··· > jk such that the smallest rectangle containing them is a subset of the polyomino.
- se(M) = the length of the largest se-chain in the filling M.

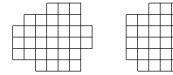


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Connections to other objects



- permutations and words
- graphs
- set partitions, linked partitions, matchings
- crossings and nestings of edges in the picture







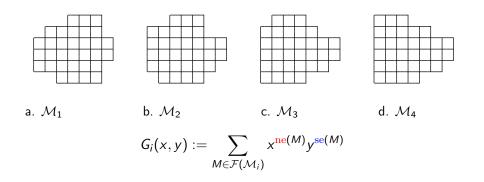


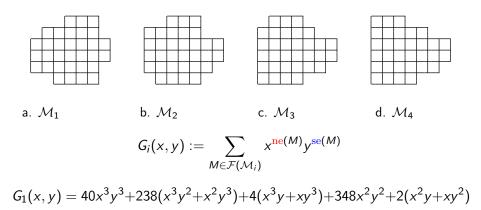
a. \mathcal{M}_1

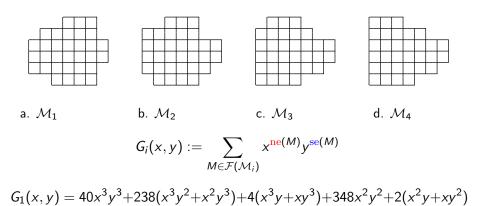
b. \mathcal{M}_2

c. \mathcal{M}_3

d. \mathcal{M}_4







$$G_2(x,y) = G_3(x,y) = G_4(x,y) = G_1(x,y)$$

Conjectures for moon polyominoes

Conjecture

For a moon polyomino \mathcal{M} , if

$$G(x,y) = \sum_{M \in \mathcal{F}(\mathcal{M})} x^{\operatorname{ne}(M)} y^{\operatorname{se}(M)}$$

then

$$G(x,y)=G(y,x).$$

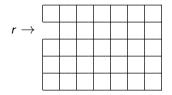
Conjecture

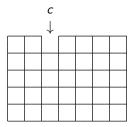
If \mathcal{M}' is obtained by permuting the rows and/or columns of $\mathcal{M},$ then

$$\sum_{M\in\mathcal{F}(\mathcal{M}')} x^{\mathbf{ne}(M)} y^{\mathbf{se}(M)} = \sum_{M\in\mathcal{F}(\mathcal{M})} x^{\mathbf{ne}(M)} y^{\mathbf{se}(M)}.$$

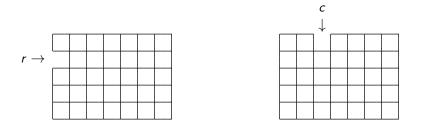
- Chen, Deng, Du, Stanley, Yan (2007) set partitions
- Backelin, West , Xin (2007) + Krattenthaler (2006) + de Mier (2006)
 Ferrers shapes
- Jonsson (2007) + Jonsson and Welker (2007) stack polyominoes
- Rubey (2011) moon polyominoes for ne-chains
- Poznanović and Yan (2014) almost moon polyominoes for ne-chains

Almost-moon: one exception to the convexity rule



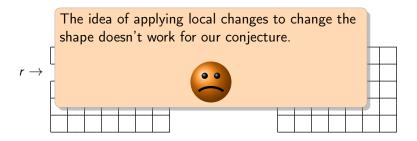


Almost-moon: one exception to the convexity rule



$$\begin{split} G_1(x,y) &= (15x^5y^3 + 13x^3y^5) + 56x^4y^4 + (80x^5y^2 + 82x^2y^5) + (1180x^4y^3 + 1178x^3y^4) \\ &\quad + 5(x^5y + xy^5) + (1210x^4y^2 + 1212x^2y^4) + 5370x^3y^3 + 10(x^4y + xy^4) \\ &\quad + (1477x^3y^2 + 1473x^2y^3) + 64x^2y^2. \\ G_2(x,y) &= (8x^5y^3 + 15x^3y^5) + 48x^4y^4 + (83x^5y^2 + 77x^2y^5) + (1129x^4y^3 + 1174x^3y^4) \\ &\quad + (9x^5y + 8xy^5) + (1273x^4y^2 + 1227x^2y^4) + 5434x^3y^3 + (6x^4y + 7xy^4) \\ &\quad + (1415x^3y^2 + 1467x^2y^3) + 60x^2y^2. \end{split}$$

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word $w \leftrightarrow$ pair of tableaux (P(w), Q(w)) of same shape

(Buch, Kresch, Shimozono, Tamvakis, Alexander Yong, 2008)

- *P*(*w*) is an increasing tableau: across rows and down columns, possibly repeated entries
- Q(w) is a <u>set-valued tableau</u>: multiple entries per cell, increasing across rows and down columns

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For w = 32412143:



What's the connection?

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Theorem (Thomas, Yong, 2011)

For a word w, let P(w) be the Hecke insertion tableau of w. Then

- lis(w) = # columns of P(w)
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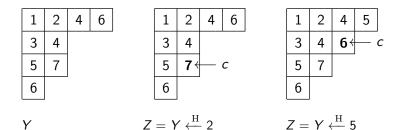
For w = 32412143:



To Hecke insert an integer x into a tableau Y:

- x is inserted in the first row R; if an output integer y is produced, then it's inserted in the next row, etc.
- If x is larger than or equal to all entries in R, then there are 2 possibilities:
 - If adding x as a new box to the first row results in an increasing tableau, do that and stop.
 - Otherwise, stop.
- Otherwise, x is strictly smaller than some entry in the first row. Let y be the smallest integer in R that is strictly larger than x.
 - If replacing y with x results in an increasing tableau, then replace y with x and insert y into the next row.
 - If replacing y with x does not result in an increasing tableau, then insert y into the next row and do not change R.

Hecke insertion: to construct P(w) and Q(w)

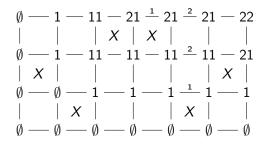


• *c* is at the bottom of the column of *Z* containing the rightmost box of the row in which the algorithm stops.

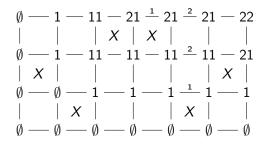
•
$$P(w) = (\cdots ((\emptyset \xleftarrow{H} w_1) \xleftarrow{H} w_2) \cdots \xleftarrow{H} w_n)$$

• Q(w) is built by inserting k in the box c that gets recorded when w_k is inserted.

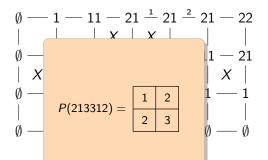
(Clemson)



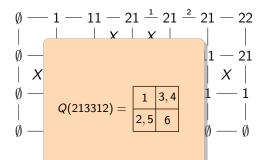
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K-jeu de taquin: $jdt_{\mathcal{C}}(\mathcal{T})$

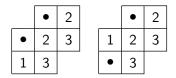
(Thomas, Yong, 2009)

- T is an increasing skew tableau
- C = set of inner corners of T (on the upper left side)
- In step k: if k is directly next to a •, swap k and •

	•	2
•	2	3
1	3	

(Thomas, Yong, 2009)

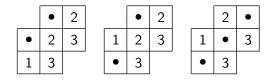
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Swap \bullet and 1

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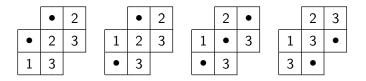
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Swap • and 2

(Thomas, Yong, 2009)

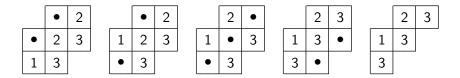
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Swap \bullet and 3

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 $\operatorname{jdt}_{C}(T)$

Definition

Two words are said to be K-Knuth equivalent if one can be obtained from the other via a finite series of applications of the following K-Knuth relations:

$$xzy \equiv zxy \qquad (x < y < z)$$

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• Words that have the same Hecke insertion tableau *P* are *K*-Knuth equivalent.

Recall

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Theorem (Thomas, Yong, '09)

If $w_1 \equiv w_2$ then $\operatorname{lis}(w_1) = \operatorname{lis}(w_2)$ and $\operatorname{lds}(w_1) = \operatorname{lds}(w_2)$.

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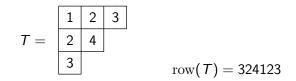
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Theorem (Buch, Samuel, '16)

Let [a, b] be an integer interval. Let $w_1 \equiv w_2$. For i = 1, 2, let $w_i|_{[a,b]}$ be the word obtained from w_i by deleting all integers not contained in the interval [a, b]. Then $w_1|_{[a,b]} \equiv w_2|_{[a,b]}$.

Properties linking all this

For an increasing tableau T, row(T) is the reading word of T, obtained by reading the entries of T from left to right along each row, starting from the bottom row and moving upward.



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$$T = \begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 \\ 3 \end{bmatrix}$$
 row(T) = 324123

Theorem (Gaetz, Mastrianni, Patrias, Peck, Robichaux, Schwein, Tam, '16) $w \equiv \operatorname{row}(P(w))$

Clemso	

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Theorem (Gaetz, Mastrianni, Patrias, Peck, Robichaux, Schwein, Tam, '16)

$$w \equiv \operatorname{row}(P(w))$$

For example,

$$P(32412143) = T \implies 32412143 \equiv 324123$$

(Clemson)

Theorem (Buch, Samuel, '16)

Let T and T' be increasing tableaux. Then $row(T) \equiv row(T')$ if and only if T and T' are K-jeu de taquin equivalent.

Theorem (Chen, Guo, Pang, '15)

Let w be a word of positive integers, and k be the maximal element appearing in w. Let w' be the word obtained from w by deleting the elements equal to k. Then P(w') is obtained from P(w) by deleting the squares occupied with k.

Theorem

If $\mathcal M$ and $\mathcal M'$ are stack polyominos with same row lengths then

 $G_{\mathcal{M}'}(x,y) = G_{\mathcal{M}}(x,y).$

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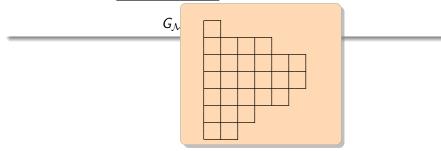
$$G_{\mathcal{M}'}(x,y) = G_{\mathcal{M}}(x,y).$$

$$\mathcal{G}_{\mathcal{M}}(x,y) = \sum_{M \in \mathcal{F}(\mathcal{M})} x^{\mathrm{ne}(M)} y^{\mathrm{se}(M)}$$



Theorem

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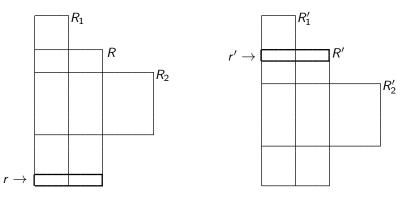


(Clemson)

Approach: Bijection between fillings of very similar shapes

Theorem

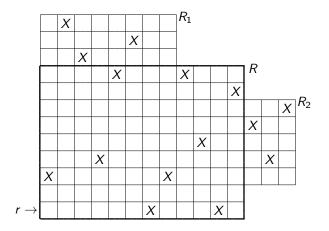
Let \mathcal{M} be a stack polyomino and let \mathcal{M}' be obtained by moving the bottom row of \mathcal{M} up. Then $G_{\mathcal{M}'}(x, y) = G_{\mathcal{M}}(x, y)$.

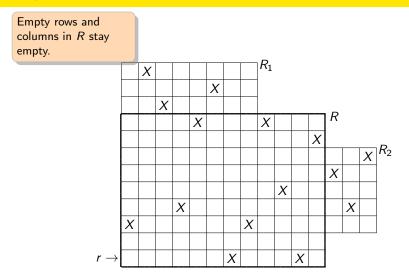


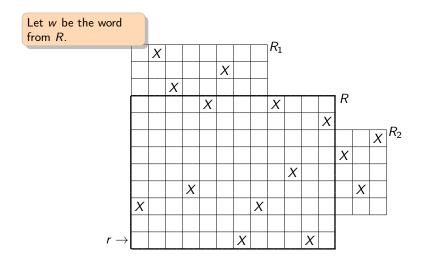
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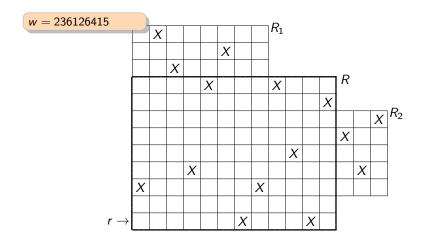


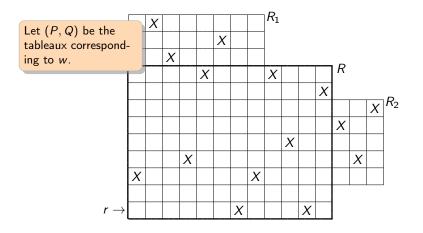
Everything outside of R stays the same.

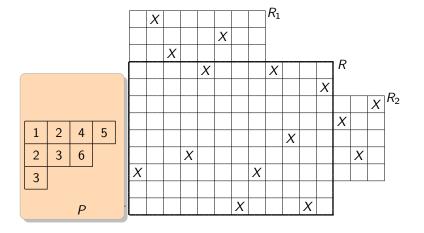


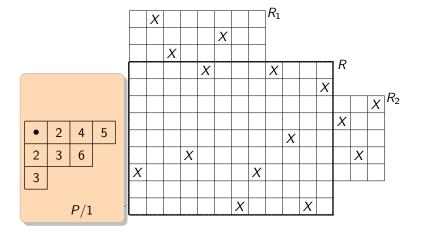


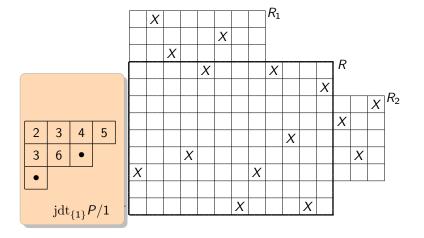


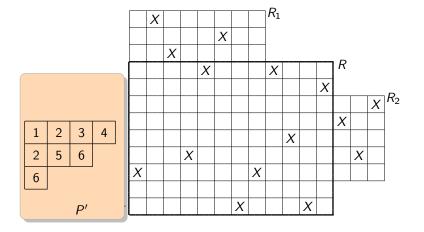


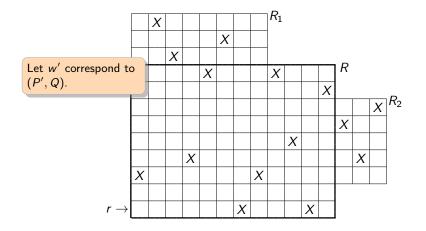


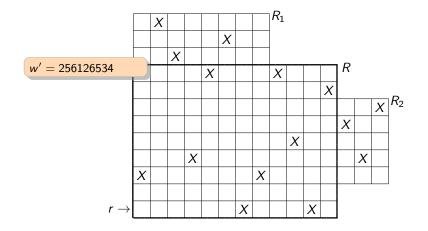


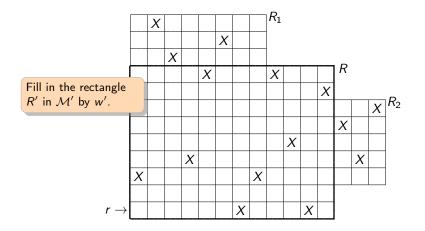




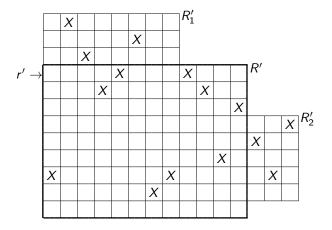






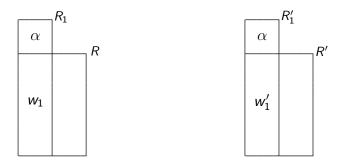


The resulting filling of \mathcal{M}'

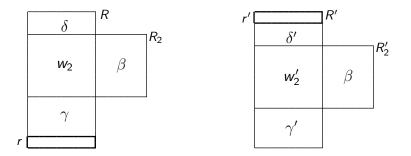


ne and se are preserved

- w and w' have Hecke insertion tableaux of same shape, so lis(w') = lis(w) and lds(w') = lds(w).
- The fillings in R_1 and R'_1 have the same recording set valued tableau.



ne and se are preserved

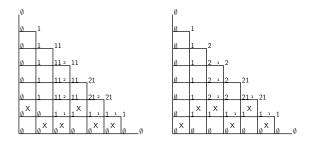


$$w \setminus 1 \equiv \operatorname{row}(P) \setminus 1 = \operatorname{row}(P/1) \equiv \operatorname{row}(P'/m) = \operatorname{row}(P') \setminus m \equiv w' \setminus m$$
$$\gamma + w_2 + \delta \equiv \gamma' + w'_2 + \delta'$$
$$w_2 \equiv w'_2$$
$$w_2\beta \equiv w'_2\beta \text{ so } \operatorname{lis}(w_2\beta) = \operatorname{lis}(w'_2\beta) \text{ and } \operatorname{lds}(w_2\beta) = \operatorname{lds}(w'_2\beta)$$

- Moving an arbitrary row up in the same way doesn't work.
- This construction doesn't preserve (ne, se) for non-stack polyominoes

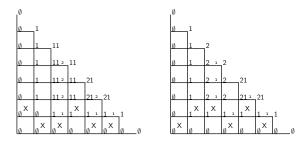
The conjecture for general moon polyominoes is still open!

Maximal crossings and nestings in linked partitions



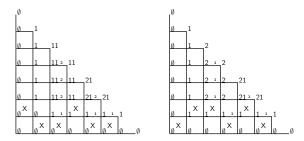
Maximal crossings and nestings in linked partitions

(Chen, Guo, Pang, '15)

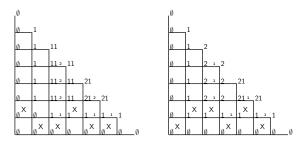


• Reflect the first tableau about the *x*-axis

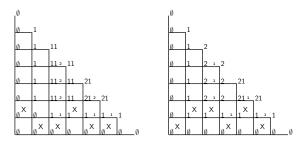
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- Apply the 'move bottom row up' rule several times



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- This produces a filling of the same initial shape with (ne, se) reversed



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- This produces a filling of the same initial shape with (ne, se) reversed
- This is a different bijection from the one described above

Thank you.